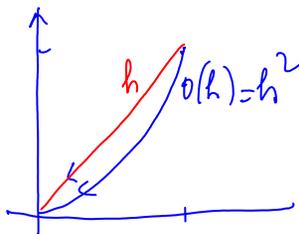


$h, o(h)?$

↳ approaches 0 faster than  $h$  does

$o(h) = h^2$

$\lim_{h \rightarrow 0} \frac{o(h)}{h} = 0$



$\frac{h^2}{h} = h$

$\Pr\{N(h) = 1\} = \lambda h + o(h)$

$\Pr\{N(h) \geq 2\} = o(h)$



Interpretation. Recall  $e^x = 1 + x + \frac{1}{2!}x^2 + \frac{1}{3!}x^3 + \dots$

$\Pr\{N(t) = n\} = e^{-\lambda t} \frac{(\lambda t)^n}{n!}, n = 0, 1, \dots$

$\Pr\{N(h) = 1\} = P_1(h) = \lambda h e^{-\lambda h} = \lambda h \left[ 1 - \lambda h + \frac{1}{2}(\lambda h)^2 + o(h) \right]$   
 $= \lambda h + o(h)$

$\Pr\{N(h) \geq 2\} = 1 - P_0 - P_1 = o(h)$

$\Pr\{N(h) = 1\} = \lambda h + o(h)$  ← Small quant  
 rate' interval width

Proof of  $\Pr\{N(t) = n\} = e^{-\lambda t} \frac{(\lambda t)^n}{n!}, n = 0, 1, 2, \dots$

Induction:  $\sum_{i=1}^n i = \frac{1}{2} n(n+1)$   
 $n=1: 1 = \frac{1}{2} \cdot 1 \cdot 2 \checkmark$

$n$	$n^2$	$3n$
1	1	3
2	4	6
3	9	9
4	16	12

rate' interval width

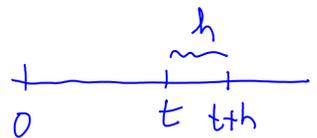
Proof of  $\Pr\{N(t)=n\} = e^{-\lambda t} \frac{(\lambda t)^n}{n!}, n=0,1,2,\dots$

$n^2 \leq 3n$		
n	$n^2$	$3n$
1	1	< 3
2	4	< 6
3	9	= 9
4	16	> 12

Induction  $\sum_{i=1}^n i = \frac{1}{2} n(n+1)$   
 $n=1: 1 = \frac{1}{2} \cdot 1 \cdot 2 \checkmark$   
 True for  $n-1$   $\sum_{i=1}^{n-1} i = \frac{1}{2} (n-1)n$   
 Show statement is true

for  $n=0: P_0(t) = e^{-\lambda t} ?$

$$P_n(t) = \Pr\{N(t)=n\} = e^{-\lambda t} \frac{(\lambda t)^n}{n!}$$



$$\begin{aligned} P_0(t+h) &= \Pr\{N(t+h)=0\} \\ &= \Pr\{N(t)=0, N(t+h)-N(t)=0\} \\ &= \Pr\{N(t)=0\} \cdot \Pr\{N(t+h)-N(t)=0\} \\ &= \Pr\{N(t)=0\} \cdot \Pr\{N(h)=0\} \\ &= P_0(t) [1 - \lambda h + o(h)] \\ &= P_0(t) - \lambda h P_0(t) + o(h) \end{aligned}$$

ind + incr

$$\lim_{h \rightarrow 0} \frac{f(t+h) - f(t)}{h} = f'(t)$$

$$\frac{P_0(t+h) - P_0(t)}{h} = \frac{-\lambda h P_0(t) + o(h)}{h}$$

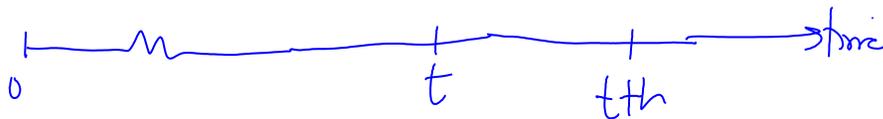
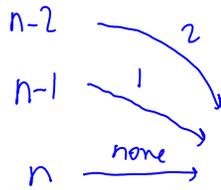
$$\lim_{h \rightarrow 0} \frac{P_0(t+h) - P_0(t)}{h} = -\lambda P_0(t) + \lim_{h \rightarrow 0} \frac{o(h)}{h}$$

$h \rightarrow 0$        $h$        $h \rightarrow 0$

$$\left. \begin{aligned} P_0'(t) &= -\lambda P_0(t) \\ P_0(0) &= 1 \end{aligned} \right\} P_0(t) = e^{-\lambda t}$$

for  $n \geq 1$ . Assume true for  $n-1$

$$Pr\{N(t) = n-1\} = e^{-\lambda t} \frac{(\lambda t)^{n-1}}{(n-1)!}$$



$$\begin{aligned} P_n(t+h) &= P_n(t) \cdot [1 - \lambda h + o(h)] + P_{n-1}(t) \cdot [\lambda h + o(h)] + P_{n+1}(t) \cdot o(h) \\ &= P_n(t) - \lambda h P_n(t) + P_{n-1}(t) \lambda h + o(h) \end{aligned}$$

$$\frac{P_n(t+h) - P_n(t)}{h} = -\lambda P_n(t) + \lambda P_{n-1}(t) + \frac{o(h)}{h}$$

In limit,

$$P_n'(t) = -\lambda P_n(t) + \lambda P_{n-1}(t)$$

$$P_n(0) = 0$$

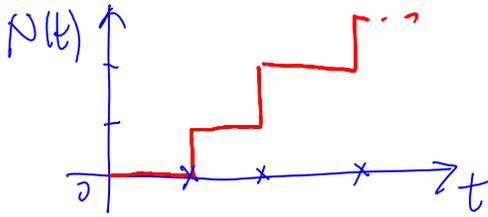
DO IT!

Using  $P_{n-1}(t)$  assumption & solving the ODE

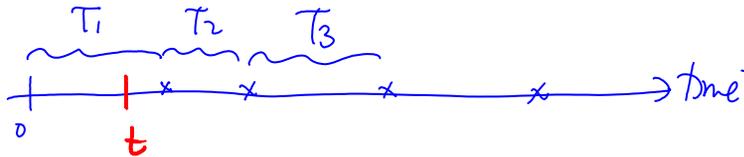
$$P_n(t) = e^{-\lambda t} \frac{(\lambda t)^n}{n!}, \quad n = 0, 1, 2, \dots$$

$$E[N(t)] = \lambda t$$

Halmos



(iii) Intercarrival times when arrivals Poisson



$T_n$ : time between  $(n-1)$ st &  $n$ th event

$$\{T_n; n \geq 1\}$$

$$T_1: f_{T_1}(t) = ? \longleftarrow F_{T_1}(t) = \Pr\{T_1 \leq t\}$$

$$\uparrow$$

$$\bar{F}_{T_1}(t) = \Pr\{T_1 > t\} \longleftarrow \text{find}$$

$$\Pr\{T_1 > t\} = \Pr\{N(t) = 0\} = e^{-\lambda t} \frac{(\lambda t)^0}{0!} = e^{-\lambda t}$$

$$\rightarrow \Pr\{T_1 \leq t\} = 1 - e^{-\lambda t} \rightarrow f_{T_1}(t) = \lambda e^{-\lambda t}$$

Similarly

$$\Pr\{T_2 > t \mid T_1 = s\} \stackrel{\text{show}}{=} e^{-\lambda t}$$

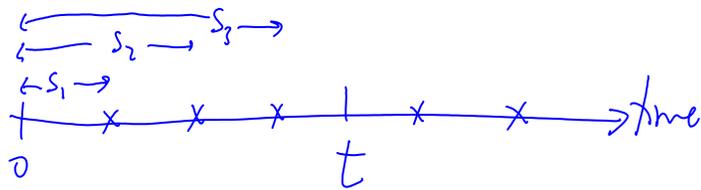
Q:  $S_n = T_1 + \dots + T_n = \sum_{i=1}^n T_i$

$$f_{S_n}(t) = \lambda e^{-\lambda t} \frac{(\lambda t)^{n-1}}{(n-1)!}, t \geq 0,$$

$\therefore \Delta$  - ...

v) An equivalent derivation of  $f_{S_n}(t)$ :

Find  $\Pr\{S_n \leq t\} \rightarrow$  diff



$$S_3 \leq t \Leftrightarrow N(t) \geq 3$$

$$S_n \leq t \Leftrightarrow N(t) \geq n$$

$$\Pr\{S_n \leq t\} = F_{S_n}(t) = \Pr\{N(t) \geq n\} = \sum_{j=n}^{\infty} e^{-\lambda t} \frac{(\lambda t)^j}{j!}$$

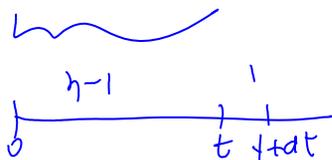
↑  
Poisson

Differentiate  $\rightarrow$  (Do!)

$$f_{S_n}(t) = \lambda e^{-\lambda t} \frac{(\lambda t)^{n-1}}{(n-1)!}, \quad n=1, 2, \dots$$

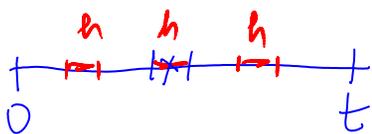
Comment. Re-write as

$$f_{S_n}(t) dt = \left[ e^{-\lambda t} \frac{(\lambda t)^{n-1}}{(n-1)!} \right] \cdot \lambda dt$$

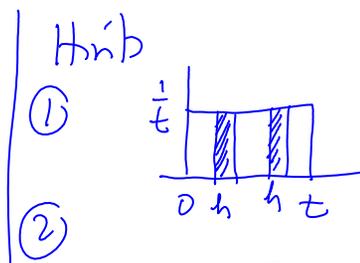


$\boxed{n^{\text{th}}}$   $n^{\text{th}}$   
 $n^{\text{th}}$   
 $n^{\text{th}}$

c) Conditional distribution of arrival times

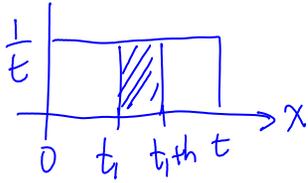


$X_i$ : time of occurrence of one



even

Theorem:  $f_{X_1|N(t)=1}(x|1) = \frac{1}{t}$ ,  $0 \leq x \leq t$  | (3)  $\Pr\{N(t)=1\} = \lambda t$



Proof. Recall  $\lim_{h \rightarrow 0} \frac{F(t+h) - F(t)}{h} = \lim_{h \rightarrow 0} \frac{\Pr\{t \leq X \leq t+h\}}{h} = f(t)$

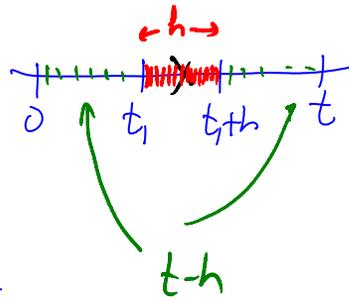
$$\Pr(A|B) = \frac{\Pr(A, B)}{\Pr(B)}$$

Now  $\Pr\{t_1 \leq X_1 \leq t_1+h | N(t)=1\}$

$$= \frac{\Pr\{t_1 \leq X_1 \leq t_1+h, N(t)=1\}}{\Pr\{N(t)=1\}}$$

$$= \frac{\Pr\{1 \text{ in } [t_1, t_1+h], \text{ zero elsewhere}\}}{\Pr\{N(t)=1\}}$$

$$= \frac{\{\lambda h e^{-\lambda h}\} \{e^{-\lambda(t-h)}\}}{\lambda t e^{-\lambda t}} = \frac{h}{t}$$



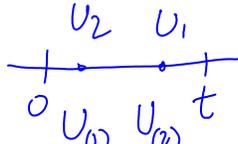
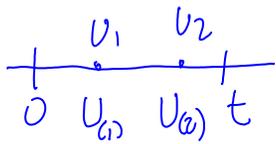
$$\therefore \lim_{h \rightarrow 0} \frac{\Pr\{ \cdot | \cdot \}}{h} = f_{X_1|N(t)=1}(x|1) = \frac{h}{t} = \frac{1}{t}, 0 \leq x \leq t$$

QED

Generalization: Preliminary result

Let  $U_1, U_2$  be uniform over  $(0, t)$

$U_{(1)}, U_{(2)}$  " order statistics for  $U_1$  and  $U_2$



$$f_{U_1}(u_1) = \frac{1}{t}$$

$$f_{U_2}(u_2) = \frac{1}{t}$$

$$U_{(1)} = \min(U_1, U_2)$$

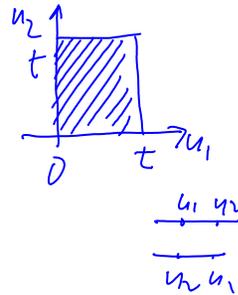
$$t=10, \quad U_1=7 \quad U_2=3$$

$$U_{(2)} = \max(U_1, U_2)$$

$$U_{(1)}=3, \quad U_{(2)}=7$$

Fact

$$f_{U_1, U_2}(u_1, u_2) = \frac{1}{t} \cdot \frac{1}{t} = \frac{1}{t^2}, \quad 0 \leq u_1 \leq t, \quad 0 \leq u_2 \leq t$$

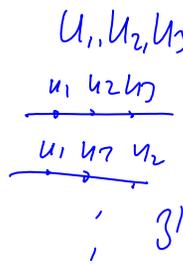
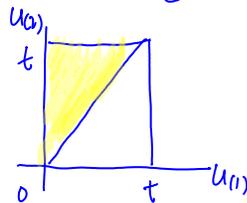


Show  $\int_0^t \int_0^t f_{U_1, U_2}(u_1, u_2) du_1 du_2 = 1$

In general  $f_{U_1, \dots, U_n}(u_1, \dots, u_n) = \frac{1}{t^n}, \quad 0 \leq u_i \leq t, \quad i=1, \dots, n$

Theorem for order statistics  $U_{(1)}$  and  $U_{(2)}$ ,

$$f_{U_{(1)}, U_{(2)}}(u_{(1)}, u_{(2)}) = \frac{2!}{t^2}, \quad 0 \leq u_{(1)} \leq u_{(2)} \leq t$$



Show  $\int_0^t \int_0^{u_{(2)}} \frac{2!}{t^2} du_{(1)} du_{(2)} = 1$

Proof. Skip  $\blacksquare$

In general  $f_{U_{(1)}, \dots, U_{(n)}}(u_{(1)}, \dots, u_{(n)}) = \frac{n!}{t^n}, \quad 0 \leq u_{(1)} \leq \dots \leq u_{(n)} \leq t$

Remark

$$\sum_{i=1}^n U_{(i)} = \sum_{i=1}^n U_i \Rightarrow$$

$$E\left(\sum_{i=1}^n U_{(i)}\right) = E\left(\sum_{i=1}^n U_i\right)$$

Theorem Given that  $N(t) = n$ , the  $n$  ordered arrival times  $S_1, S_2, \dots, S_n$  have the same distribution as the order statistics corresponding to the  $n$  i.i.d. uniform r.v.'s over  $(0, t)$ , i.e.,

$$f_{S_1, \dots, S_n | N(t)=n}(s_1, \dots, s_n | t) = \frac{n!}{t^n}, \quad 0 \leq s_1 \leq \dots \leq s_n \leq t$$

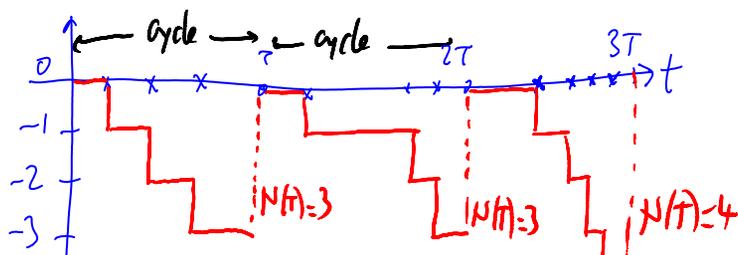
Ex. Periodic-review inventory policy

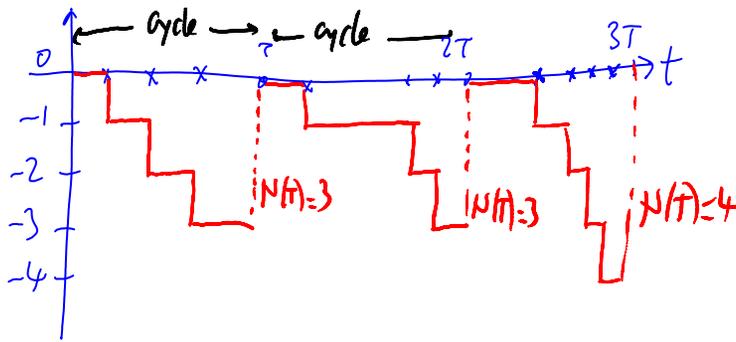
Renew inventory at  $T, 2T, 3T, \dots$  ( $T$ : review length)

Suppose all demand is backordered (backlogged) & order is placed at end of period to satisfy all demands.

Demand Process with rate  $\lambda$

Unit purchase cost =  $c$





## Renewal reward Thm (RRT)

$$AC(T) = C(T) = \frac{E(\text{cycle cost})}{E(\text{cycle length})}$$

• Cycle length =  $T$  (time)

$$\lambda : \left( \frac{\text{units}}{\text{time}} \right)$$

$$c : \left( \frac{\$}{\text{unit}} \right)$$

$$E(N(T)) = \lambda T$$

• Cycle cost

(i) Order cost :  $K$  [ $\$$ ]

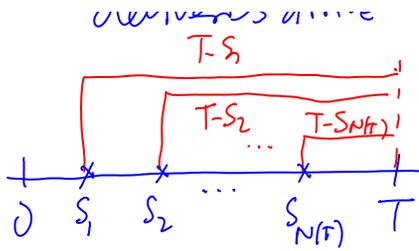
(ii) Purchase cost :  $cN(T) \Rightarrow E(\text{Purchase cost}) = c\lambda T$

$$\cancel{\frac{\$}{u}} \cdot \frac{u}{\cancel{t}} \neq$$

(iii) Backorder cost

$$\hat{w} : \text{cost of backorder/unit/time} \quad \frac{\left[ \frac{\$}{u} \right]}{\left( \frac{t}{i} \right)} = \left[ \frac{\$}{u} \cdot \frac{1}{t} \right]$$

$W(T)$  : total waiting time of all customers (units)  
deliveries arrive



$$\begin{aligned}
 W(T) &= (T - s_1) + (T - s_2) + \dots + (T - s_{N(F)}) \\
 &= \sum_{i=1}^{N(F)} (T - s_i)
 \end{aligned}$$